Application of Factorial and Binomial Identities in Information, Cybersecurity and Machine Learning

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-----ABSTRACT-----

This paper presents application of the binomial and factorial identities and expansions that are used in artificial intelligence, machine learning, and cybersecurity. The factorial and binomial identities can be used as methodological advances for various algorithms and applications in information and computational science. Cybersecurity is the practice of protecting the computing systems, communication networks, data and programs from cyber-attacks. Its objective is to reduce the risk of cyber-attacks and protect against the unauthorized exploitation of systems and networks. For this purposes, we need a strong cryptographic algorithms like RSA algorithm and Elliptic Curve Cryptography. In this connection, computing and combinatorial techniques based on factorials and binomial distributions are developed for the researchers who are working in artificial intelligence and cybersecurity.

Keywords - Cybersecurity, combinatorics, computation, factorial, information.

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 I. INTRODUCTION
 Then, (3 + 1)(3 + 2)(3 + 3)(3 + 4)(3 + 5)

 Computational science is a rapidly growing multi-and
 Then, (3 + 1)(3 + 2)(3 + 3)(3 + 4)(3 + 5)

 $= 6720 = 56 \times 120 = 56 \times 5!$ and
 $(5 + 1)(5 + 2)(5 + 3) = 336 = 56 \times 6.$

inter-disciplinary area where science, engineering, computation, mathematics, and collaboration uses advance computing capabilities to understand and solve the most complex real life problems. Wireless communication is not fully secure for transmission of information in a peerto-peer network, that is, some information leakage through the transmission of wireless signal is unavoidable. In this case, cybersecurity is the practice of protecting the computing systems, devices, communication networks, programs and data from cyber-attacks. The aim of this article is to reduce the risk of cyber-attacks and protect against the unauthorized exploitation of systems and networks. For this purposes, we need a strong security mathematical algorithm like RSA algorithm and Elliptic Curve Cryptography. The factorial [1, 2] and binomial theorem [3-7] will help to build a strong artificial intelligence-based cryptographic algorithm.

II. THEOREM IN FACTORIALS

The factorial of a non-negative integer *n*, denoted by *n*!, is the product of all positive integers less than or equal to *n*. For example, $4! = 1 \times 2 \times 3 \times 4 = 24$ and 0! = 1.

Theorem in Factorials $[11] : (n_1 + n_2 + n_3 + \dots + n_k)! = T \times n_1! \times n_2! \times n_3! \times \dots \times n_k!$, where $T, n_i \in N = \{1, 2, 3, \dots\} \& i = 1, 2, 3, \dots, k$.

Proof: Let $x = n_2 + n_3 + \dots + n_k$. $(n_1 + x)! = n_1! \times (n_1 + 1)(n_1 + 2)(n_1 + 3) \dots (n_1 + x)$. We know that $(r + 1)(r + 2)(r + 3) \dots (r + n)$ $= a \times n!$, where *a* is a positive integer. For example, Let n = 5 and r = 3. From the above result, we get $n_1! \times (n_1 + 1)(n_1 + 2)$ $(n_1 + 3) \cdots (n_1 + x) = a_1 \times n_1! \times x!$, *i.e.*, $(n_1 + x)! = a_1 \times n_1! \times x! = a_1 \times n_1!$ $\times (n_2 + n_3 + \cdots + n_k)! (\because x = n_2 + n_3 + \cdots + n_k)$.

Similarly, if we continue the same process up to *k*-1 times, then we obtain the result: $(n_1 + n_2 + n_3 + \dots + n_k)! = (a_1 \times a_2 \times a_3 \times \dots \times a_{k-1}) \times n_1! \times n_2! \times n_3! \times \dots \times n_k!$.

Let $T = (a_1 \times a_2 \times a_3 \times \dots \times a_{k-1})$, where $T, a_i \in N = \{1, 2, 3, \dots \} \& i = 1, 2, 3, \dots, k-1$. Then, $(n_1 + n_2 + n_3 + \dots + n_k)!$ $= T \times n_1! \times n_2! \times n_3! \times \dots \times n_k!$.

Hence, theorem is proved.

III. ARTIFICIAL INTELLIGENCE-BASED CYBERSECURITY

The mathematical results mentioned-below can be used as an application in the artificial intelligence-based cybersecurity and these results also help to create a strong security algorithm like RSA algorithm and Elliptic Curve Cryptography in the field of cybersecurity in order toreduce the risk of cyber-attacks and protect against the unauthorized exploitation of systems and networks.

We have already understood the below factorial identity and detailed proof in the previous section of this paper. Now, we find how to use this result in another way.

$$(n_1 + n_2 + n_3 + \dots + n_k)! = (a_1 \times a_2 \times a_3 \times \dots \times a_{k-1})$$
$$\times n_1! \times n_2! \times n_3! \times \dots \times n_k!,$$

that is, $(n_1 + n_2 + n_3 + \dots + n_k)!$ = $T \times n_1! \times n_2! \times n_3! \times \dots \times n_k!$. Here, $n_1, n_2, n_3, \dots, n_k$ are possitive integers.

An alternative way for finding the positive integer T is given below:

$$T = \frac{(n_1 + n_2 + n_3 + \dots + n_k)!}{n_1! \times n_2! \times n_3! \times \dots \times n_k!}$$

where $T, n_i \in N = \{1, 2, 3, \dots\} \& i = 1, 2, 3, \dots$

For our convenience, we can rearrange the positive integers that are equal to the product T. Let T = 64. $T = 1 \times 64$; $T = 2 \times 32$;

Let T = 04. $T = 1 \times 04$, $T = 2 \times 32$, $T = 2 \times 2 \times 16$; $T = 2 \times 4 \times 8$; $T = 4 \times 4 \times 4$; etc. Similarly, if T = 21, Then, $T = 1 \times 21$ or $T = 3 \times 7$.

3.1 Binomial Identities

Binomial identities [3-13] mentioned below can be used in cybersecurity:

$$(1) V_{r}^{n} = V_{r}^{n}(n, r \ge 1). (2)V_{r}^{n+1} - V_{r}^{n} = V_{r-1}^{n}.$$

$$(3)1 + V_{1}^{1} + V_{1}^{2} + V_{1}^{3} \cdots V_{1}^{n} = V_{2}^{n}. (4)V_{n}^{n} = 2V_{n-1}^{n}.$$

$$(5) V_{0}^{n} + V_{1}^{n} + V_{2}^{n} + V_{3}^{n} \cdots + V_{r-1}^{n} + V_{r}^{n} = V_{r}^{n+1}.$$

$$(6) \sum_{i=0}^{n} (i+1)V_{i}^{n-i} = (n+2)2^{n-1}.$$

$$(7) \sum_{i=0}^{n} i \times V_{i}^{n-i} = n2^{n-1}. (8) \sum_{i=0}^{n} V_{i}^{n-i} = 2^{n}.$$

$$(9) \sum_{i=0}^{n} (i-1)V_{i}^{n-i} = (n-2)2^{n-1}.$$

$$(10) V_{1}^{1} + V_{2}^{2} + V_{3}^{3} + \cdots + V_{n}^{n} = 2(V_{0}^{1} + V_{1}^{2} + V_{2}^{3} + \cdots + V_{n-1}^{n}).$$

$$(11) \sum_{i=1}^{r} V_{i}^{n+1} = \sum_{i=0}^{r} V_{i}^{0} + \sum_{i=0}^{r} V_{i}^{1} + \sum_{i=0}^{r} V_{i}^{2} + \sum_{i=0}^{r} V_{i}^{3} + \cdots + \sum_{i=r}^{r} V_{i}^{n}$$

$$(12) \sum_{i=0}^{r} V_{i}^{n+1} x^{i} = \sum_{i=0}^{r} V_{i}^{n} x^{i} + \sum_{i=1}^{r} V_{i-1}^{n} x^{i} + \sum_{i=r}^{r} V_{i-r}^{n} x^{i}.$$

The numerical expression of binomial coefficient used in binomial identities is given below:

$$V_r^n = \frac{(r+1)(r+2)(r+3)\cdots(r+n-1)(r+n)}{n!},$$

(n,r \in N, n \ge 1, & r \ge 0).

3.2 Computation of Sum of Binomial Coefficients The computation of sum of binomial coefficients is developed by sing the binomial identity (10).

$$\begin{split} \prod_{i=1}^{1} \frac{(1+i)}{1!} + \prod_{i=1}^{2} \frac{(2+i)}{2!} + \prod_{i=1}^{3} \frac{(3+i)}{3!} + \cdots \\ &+ \prod_{i=1}^{n-1} \frac{(n-1+i)}{(n-1)!} + \prod_{i=1}^{n} \frac{(n+i)}{n!} = \\ 2 \left(\prod_{i=1}^{1} \frac{(0+i)}{1!} + \prod_{i=1}^{2} \frac{(1+i)}{2!} + \prod_{i=1}^{3} \frac{(2+i)}{3!} + \cdots + \cdots \right) \\ &+ \prod_{i=1}^{n-1} \frac{(n-2+i)}{(n-1)!} + \prod_{i=1}^{n} \frac{(n-1+i)}{n!} \right) \\ &\implies \sum_{j=1}^{n} \prod_{i=1}^{j} \frac{(j+i)}{j!} = 2 \sum_{j=1}^{n} \prod_{i=1}^{j} \frac{(j-i+i)}{j!}. \end{split}$$

3.3 Relation between the Binomial Expansions with multiple of 2

Relation 1:
$$\sum_{i=0}^{n} (i+1)V_i^{n-i} + \sum_{i=0}^{n} (i-1)V_i^{n-i}$$
$$= \sum_{i=0}^{n} i \times V_i^{n-i} = n2^{n-1}.$$

Proof: Let us simply the general terms in the two parts of binomial expansions (Relation 1) as follows:

$$\begin{aligned} &(i+1)V_{i}^{n-i} + (i-1)V_{i}^{n-i} \\ &= 2iV_{i}^{n-i}. \text{ This idea can be applied for Relation 1.} \\ &\sum_{i=0}^{n} (i+1)V_{i}^{n-i} + \sum_{i=0}^{n} (i-1)V_{i}^{n-i} = 2\sum_{i=0}^{n} iV_{i}^{n-i} \\ &= (n+2)2^{n-1} + (n-2)2^{n-1} \\ &= 2n2^{n-1}. \end{aligned}$$

$$\begin{aligned} &\text{Then, } 2\sum_{i=0}^{n} iV_{i}^{n-i} = 2n2^{n-1} \Longrightarrow \sum_{i=0}^{n} iV_{i}^{n-i} = n2^{n-1}. \end{aligned}$$

$$\begin{aligned} &\text{Relation } 2: \sum_{i=0}^{n} (i+1)V_{i}^{n-i} - \sum_{i=0}^{n} (i-1)V_{i}^{n-i} = \sum_{i=0}^{n} V_{i}^{n-i} \\ &= 2^{n} \end{aligned}$$

Proof: Let us simply the general terms in the two parts of binomial expansions (Relation 2) as follows: $(i + 1)V^{n-i} - (i - 1)V^{n-i}$

$$= 2V_{i}^{n-i}.$$
 This idea can be applied for Relation 2.

$$\sum_{i=0}^{n} (i+1)V_{i}^{n-i} - \sum_{i=0}^{n} (i-1)V_{i}^{n-i} = 2\sum_{i=0}^{n} V_{i}^{n-i}$$

$$= (n+2)2^{n-1} - (n-2)2^{n-1}$$

$$= 4 \times 2^{n-1}.$$
Then, $2\sum_{i=0}^{n} V_{i}^{n-i} = 22^{n} \Longrightarrow \sum_{i=0}^{n} V_{i}^{n-i} = 2^{n}.$
Hence, two relations are proved

Hence, two relations are proved.

3.4 Computation of Binomial Expansions Computation of binomial expansions [6-8] is constituted by using the binomial identity (12),

$$\begin{split} \sum_{i=1}^{r} V_i^{n+1} &= \sum_{i=0}^{r} V_i^0 + \sum_{i=0}^{r} V_i^1 + \sum_{i=0}^{r} V_i^2 + \sum_{i=0}^{r} V_i^3 + \cdots \\ &+ \sum_{i=0}^{r} V_i^{n-1} + \sum_{i=0}^{r} V_i^n \implies \end{split}$$

$$\begin{split} \sum_{i=1}^{r} \frac{(i+1)(i+2)(i+3)\cdots(i+n)}{n!} \\ &= \sum_{i=0}^{r} \frac{1}{0!} + \sum_{i=0}^{r} \frac{(i+1)}{1!} \\ &+ \sum_{i=0}^{r} \frac{(i+1)(1+2)}{2!} + \cdots \\ &+ \sum_{i=0}^{r} \frac{(i+1)(i+2)(i+3)}{3!} + \cdots \\ &+ \sum_{i=0}^{r} \frac{(i+1)(i+2)(i+3)\cdots(i+n-1)}{(n-1)!} . \end{split}$$

3.5 Summation of Binomial Series

The following binomial series [7] is constituted based on the binomial identity (13).

$$\sum_{i=0}^{r} V_{i}^{n+1} x^{i} = \sum_{i=0}^{r} V_{i}^{n} x^{i} + \sum_{i=1}^{r} V_{i-1}^{n} x^{i} + \sum_{i=2}^{r} V_{i-2}^{n} x^{i} + \cdots + \sum_{i=r-1}^{r} V_{i-(r-1)}^{n} x^{i} + \sum_{i=r}^{r} V_{i-r}^{n} x^{i}.$$

Proof: Let's show that the computation of addition of binomial series (right-hand side) is equal to the sum of binomial series for upper limit r+1 (left-hand side).

$$\sum_{i=0}^{r} V_{i}^{n+1} x^{i} = \sum_{i=0}^{r} V_{i}^{n} x^{i} + \sum_{i=1}^{r} V_{i-1}^{n} x^{i} + \sum_{i=2}^{r} V_{i-2}^{n} x^{i} + \cdots + \sum_{i=r-1}^{r} V_{i-(r-1)}^{n} x^{i} + \sum_{i=r}^{r} V_{i-r}^{n} x^{i}$$
$$= (V_{0}^{n} + V_{1}^{n} x + V_{2}^{n} x^{2} + V_{2}^{n} x^{3} + \cdots + V_{r}^{n} x^{r})$$

$$= (V_0^n + V_1^n x^{-1} + V_2^n x^{-1} + V_3^n x^{-1} + V_r^n x^{-1}) + (V_0^n x + V_1^n x^2 + V_2^n x^3 + V_3^n x^4 + \cdots + V_{r-1}^n x^r) + (V_0^n x^2 + V_1^n x^3 + V_2^n x^4 + V_3^n x^5 + \cdots + V_{r-2}^n x^r) + \cdots + (V_0^n x^{r-1} + V_1^n x^r) + V_0^n x^r = V_0^n + (V_0^n + V_1^n) x + (V_0^n + V_1^n + V_2^n) x^2 + \cdots$$

 $(Note that V_0^p + V_1^r + V_2^r + V_3^r + \dots + V_n^r)x^r$ (Note that $V_0^p + V_1^p + V_2^p + \dots + V_r^p = V_r^{p+1} and V_0^p$ $= V_0^{p+1} = 1$) $- V_0^{n+1} + V_0^{n+1}x^2 + V_0^{n+1}x^3 + V_0^{n+1}x^4 + \dots$

$$V_{0} + V_{1} + V_{2} + V_{2} + V_{3} + V_{3} + V_{4} + X + \cdots$$
$$+ V_{r-1}^{n+1} x^{r-1} + V_{r}^{n+1} x^{r} = \sum_{i=0}^{r} V_{i}^{n+1} x^{i}.$$

Hence, it is proved.

IV. CONCLUSION

In this article, a mathematical techniques and applications [15] for an artificial intelligence-based cybersecurity have been introduced in order to protect the computing systems, devices, networks, programs and data from cyber-attacks, that is, to reduce the risk of cyber-attacks and protect against the unauthorized exploitation of systems, programs, and networks. The factorial and binomial identities and expansions can be used as artificial intelligence-based methodological advance in cybersecurity.

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