Cost Optimization Technique of Task Allocation in Heterogeneous Distributed Computing System

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ABSTRACT

A Distributed Computing System (DCS) is a network of workstations, personal computer and/or other computing systems. Such system may be heterogeneous in the sense that the computing nodes may have different speeds and memory capacities. A DCS accepts tasks from users and executes different modules of these tasks on various nodes of the system. Task allocation in a DCS is a common problem and a good number of task allocation algorithms have been proposed in the literature. In such environment an application runs in a DCS can be accessible on every node present within the DCS. In such cases if number of tasks is less than or equal to available processors in the DCS, we can assign these task without any trouble. But this allocation becomes complicated when numbers of tasks are greater than the numbers of processors. The problem of task allocation for processing of 'm' tasks to 'n' processors (m>n) in a DCS is addressed here through a new modified tasks allocation technique. The model, presented in this paper allocates the tasks to the processor of different processing capacity to increase the performance of the DCS. The technique presented in this paper is based on the consideration of processing cost of the task to the processors. We have tried a new technique to assign all the tasks as per the required availability of processors and their processing capacity so that all the tasks of application get execute in the DCS.

Keywords - Distributed Computing System, Optimization Technique, Performance, Processing Cost, Task allocation.

1. INTRODUCTION

Distributed Computing System (DCS) aims at achieving higher execution speed than the one obtainable with uniprocessor system by exploiting the collaboration of multiple computing nodes interconnected in some fashion. The best possible speed up will obviously be obtained if the various partitions of the given computational task can run independently in parallel. In DCS a single large problem is broken into multiple small problems and processed the smaller part and combines the solution into one solution for the problem. The purpose of the DCS is to coordinate the use of shared resources or provide communication services to the users. In DCS several set of processors handles multiple tasks and ensure the execution of all the requested tasks modules. A DCS can be seen as virtual uniprocessor. In the present research paper, we have taken a domain of DCS where there are 'n' heterogeneous processors and 'm' tasks (where m>n). This scenario looks like a FIFO (First in First out) ordered queue. When first most tasks are assigned to the processor, remaining tasks have to wait until first task has been completed. As given in figure 1.

As per the presented scenario initially all the available processor will be assigned with one task for each processor and remaining tasks will be on hold until the allocated processor will be free. To solve such task allocation problem in DCS, will have to make such arrangement where more than one task can be assigned to a single processor in order to achieve minimum execution cost. Some of the task allocation schemes have been reported in the literature, such as Optimal execution cost [1,2,5,7,8,9,12,14], Task Allocation Model [6,10,11,20], Task Allocation for maximum reliability [15,16,17,19], Task Scheduling [4,13], Decentralized management of bi-modal network [3] and Disaster recovery mechanism using widely distributed
networking [18]. In this research paper we have taken the problem of task allocation in DCS based on execution cost. Here we have proposed a new task allocation scheme to get minimum execution cost by allocating task. And overall allocation should also be balanced that avoid the situation of overloading to make the proper utilization of processors of the DCS.

2. NOTATIONS

\begin{align*}
 p & \quad \text{Processor} \\
t & \quad \text{Task} \\
n & \quad \text{Number of Processors} \\
m & \quad \text{Number of Tasks} \\
TCM & \quad \text{Task Cost Matrix} \\
PCM & \quad \text{Processing Task Cost Matrix} \\
TPCM & \quad \text{Temporary Processing Cost Matrix} \\
STPCM & \quad \text{Sum of Temporary Processing Cost Matrix}
\end{align*}

3. OBJECTIVE

The objective of this research is to minimize the overall processing cost for a Distributed Computing System (DCS) through optimally assigning the tasks of various processors in heterogeneous environment, so that the performance of the DCS is to be enhanced. Processing of task is referred here as task along with all its sub tasks are executed. Assignment type would be static in DCS that will also insure the processing of all the modules of task as task modules are more than the numbers of processors of the environment. Here the performance is measured in term of processing cost of the task that have to be get processed on the processors of the environment and it have to be efficiently processed i.e., cost to be minimized.

4. TECHNIQUE

In order to evaluate the overall optimal processing cost of a DCS, we have consider the problem where a set \( P = \{p_1, p_2, p_3, \ldots, p_n\} \) of ‘n’ processors with different configuration and a set \( T = \{t_1, t_2, t_3, \ldots, t_m\} \) of ‘m’ tasks, where \( m > n \). processing cost are known for all task modules for every processor and will be arrange the processing cost for each task for different processor in a Processing Task Cost Matrix (PCM) of order \((n \times m)\). After arranging the task modules we will break PCM Matrix into another frame called Temporary Processing Cost Matrix TPCM \([3][3]\) and will also initialize the processing load for all processors by zero. In the next step, will form all possible combination of elements column wise by using TPCM (.). And will find the total number of possible combination 27 and will also calculate the sum of all possible combination and store in another matrix Sum of Temporary Processing Cost Matrix STPCM(.). After calculate the sum of all possible combination in the matrix will find minimum value of sum and will compare processing load on within the available processor if task match with the appropriate processor where the processing load is also minimum then assign the task otherwise search for next minimum value of sum and the task will get assign, as per the allocation scheme these steps will be repeated until all the task module will assign to the processors in DCS.

5. ALGORITHM

1. Start Algorithm
2. Read the number of task in m
3. Read the number of processor in n
4. Store task and Processing Cost into Matrix PCM(,) \( n \times m \) of order
5. While (All task != Assigned)
   
   \begin{itemize}
   
   \item Consider n tasks in sequence and store in TPCM(.) in order of \( n \times m \).
   \item Create all possible combination and calculate the sum for each combination and store values in STPCM().
   \item Search the minimum value of sum and draw Combination Cost Matrix (CCM)
   \item Check the processing load of matching processor for eligible task (with minimum processing cost).
   \item If it is minimum then the task will get assign otherwise will go to step iii and search next minimum value to task get assigned.
   \end{itemize}

6. State the results
7. End of algorithm

6. IMPLEMENTATION

Here in this research paper, we have taken Distributed Computing System (DCS) which consist a set \( P \) of 3 processors \( \{p_1, p_2, p_3\} \) with different configuration, and a set \( T \) of 6 tasks \( \{t_1, t_2, t_3, t_4, t_5, t_6\} \). Two imaginary tasks are also added to complete the frame with 0 processing cost \( \{t_7, t_8\} \). It is shown in the Table 1. The processing cost of each task varies processor vise and known and mentioned in the processing cost matrix namely PCM of order \( 3 \times 8 \).

\begin{table}[h]
\centering
\caption{Processing Cost Matrix}
\begin{tabular}{cccc}
\hline
 & \( t_1 \) & \( t_2 \) & \( t_3 \) & \( t_4 \) \\
\hline
\( P_1 \) & 1.522 & 1.256 & 2.025 & 1.523 \\
\( P_2 \) & 2.584 & 1.866 & 2.015 & 0.985 \\
\( P_3 \) & 0.955 & 1.658 & 2.001 & 1.221 \\
\hline
\end{tabular}
\end{table}

To accomplish task allocation process in optimize way, we take first three task to complete first frame and will create a temporary matrix TPCM of \([3][3]\) as follows:
Table 2: Temporary Processing Cost Matrix

\[
\text{TPCM} = \begin{bmatrix}
 t_1 & t_2 & t_3 \\
 P_1 & 1.522 & 1.256 & 2.025 \\
 P_2 & 2.584 & 1.866 & 2.015 \\
 P_3 & 0.955 & 1.658 & 2.001 \\
\end{bmatrix}
\]

Here we will make all the possible combinations of elements and find out the smallest sum of possible combinations.

For the first frame the combination, we have as follows:

Table 3: Sum of Temporary Cost Matrix

\[
\begin{array}{ccc}
 & t_1 & t_2 & t_3 \\
 P_1 & - & 1.256 & - \\
 P_2 & - & 1.866 & - \\
 P_3 & 0.955 & - & - \\
\end{array}
\]

Here \( t_1 \) will get assign to \( p_3 \) processor. Allocation table 1 will be as follows for the first frame:

**Allocation Table 1**

<table>
<thead>
<tr>
<th>Processor</th>
<th>Task</th>
<th>Processing Cost</th>
</tr>
</thead>
<tbody>
<tr>
<td>( P_3 )</td>
<td>( t_1 )</td>
<td>0.955</td>
</tr>
</tbody>
</table>

By the following the same procedure for the rest frames, we find following allocation table:

**Allocation Table 2**

<table>
<thead>
<tr>
<th>Processor</th>
<th>Task</th>
<th>Processing Cost</th>
</tr>
</thead>
<tbody>
<tr>
<td>( P_1 )</td>
<td>( t_2 ) * ( t_3 )</td>
<td>( 1.256 + 1.352 = 2.608 )</td>
</tr>
<tr>
<td>( P_2 )</td>
<td>( t_3 ) * ( t_4 )</td>
<td>( 2.015 + 0.985 = 3.000 )</td>
</tr>
<tr>
<td>( P_3 )</td>
<td>( t_1 ) * ( t_6 )</td>
<td>( 0.955 + 2.009 = 2.964 )</td>
</tr>
<tr>
<td>Total Processing Cost</td>
<td></td>
<td>8.572</td>
</tr>
</tbody>
</table>

7. CONCLUSION

In this research paper we have taken a problem where \( m \) number of tasks needs to assign \( n \) number of processor where \( m \) is always greater than \( n \). We did task allocation in optimize way in terms of processing cost in a Distributed Computing System (DCS). The technique introduced in research paper applied on several sets of input data and that verified the objective of get minimum processing cost for given tasks for their execution. Here we measured the performance in terms of processing cost of the tasks that has been processed by the processor of the network. The outcome of the given example as mentioned below here:

Resultant Allocation Table

<table>
<thead>
<tr>
<th>Processor</th>
<th>Task</th>
<th>Processing Cost</th>
</tr>
</thead>
<tbody>
<tr>
<td>( P_1 )</td>
<td>( t_2 ) * ( t_5 )</td>
<td>2.608</td>
</tr>
<tr>
<td>( P_2 )</td>
<td>( t_3 ) * ( t_4 )</td>
<td>3.000</td>
</tr>
<tr>
<td>( P_3 )</td>
<td>( t_1 ) * ( t_6 )</td>
<td>2.964</td>
</tr>
<tr>
<td>Total Processing Cost</td>
<td></td>
<td>8.572</td>
</tr>
</tbody>
</table>

Figure 2 shows the final task allocation as mentioned in allocation table.

**REFERENCES**


